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#### Selection of Spectral-Spatial Channels in SDM Flexgrid Optical Networks

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#### Agenda

- Introduction and Motivation
- Optimization Problem
- Algorithm
- Results
- Conclusions



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#### Internet Capacity Crunch – Why???

#### Bandwidth-hungry applications/services:

- HDTV, video streaming, 4K
- Big data processing
- Game streaming

#### • Increasing number of users/devices:

- Internet reaches almost every person on Earth
- Every user uses many devices (smartphone, iPad, PC, TV, etc)
- Internet of Things (IoT) the number of devices connected to the Internet will grow from 5 billion now up to 50 billion in 2020
- Evolution access network technologies:
  - FTTx
  - LTE 300 Mbps
  - 5G 10Gbps



#### **Cisco Traffic Forecasts**

- The **Cisco Global Cloud Index (GCI)** forecasts data center and cloud traffic and related trends
- The Cisco Visual Networking Index (VNI) is the company's ongoing effort to forecast and analyze the growth and use of IP networks worldwide
- **CAGR** (Compound Annual Growth Rate)





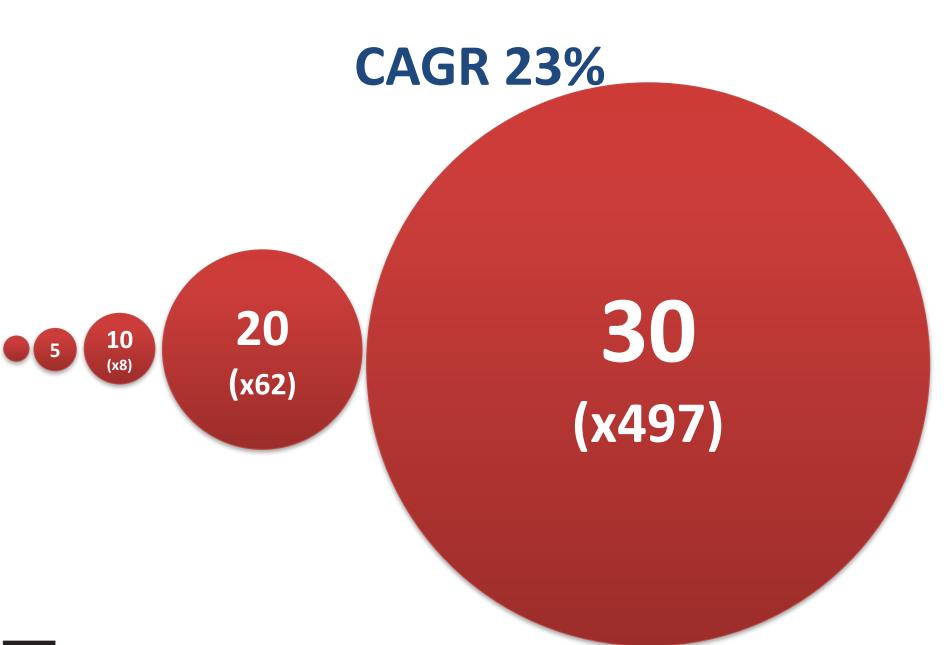


#### **Predicted CAGR**

#### **IP Traffic**

- 2013 VNI report for years 2012-2017 report, CAGR=23%
- 2014 VNI report for years 2013-2018 report, CAGR=21%
- 2015 VNI report for years 2014-2019 report, CAGR=23%
- 2016 VNI report for years 2015-2020 report, CAGR=22%
- Content Delivery Network (CDN) Traffic
- 2013 VNI report for years 2012-2017 report, CAGR=34%
- 2014 VNI report for years 2013-2018 report, CAGR=34%
- 2015 VNI report for years 2014-2019 report, CAGR=38%
- 2016 VNI report for years 2015-2020 report, CAGR=34%





S.

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#### **CAGR 30%**



# **30** (x2619)



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K. Walkowiak, "Selection of Spectral-Spatial Channels in

20

(x190)

# How to overcome the capacity crunch???

- Deliver network traffic in a smart way (CDN, anycasting, multicasting, etc.)
- Limit network traffic (blocking P2P traffic, throttling video traffic, etc.)
- Update backbone (optical) networks



## **Evolution of Optical Networks**

- **Currently**, most of the **transport optical networks use WDM** (Wavelength Division Multiplexing) technology with fixed-grid
- Possible ways to increase capacity of optical networks:
  - Elastic Optical Networks (EONs) with higher flexibility in the spectrum domain (flex-grid)
  - Space-Division Multiplexing (SDM) with higher flexibility in the space domain



#### SDM

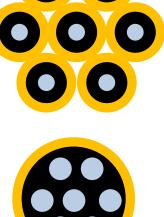
- The key idea behind SDM is to use the space domain, in which the spatial resources can be flexibly assigned to different traffic demands
- SDM allows to increase the overall transmission capacity in a cost-effective manner by integrating to a certain extent multiple transmission systems in parallel



## **SDM Technologies**

- Fiber bundle standard fibers, often deployed in bundles (to offset the costs of digging trenches)
- Multicore fiber fibers with multiple cores within a single fiber cladding, forming multicore fibers (MCFs), offer an increase in available bandwidth equal to their core count
- Multimode fiber fibers with a single, large core, which can carry additional optically-guided spatial modes, few-mode fibers (FMFs) offer a potential capacity multiplier equal to the mode count







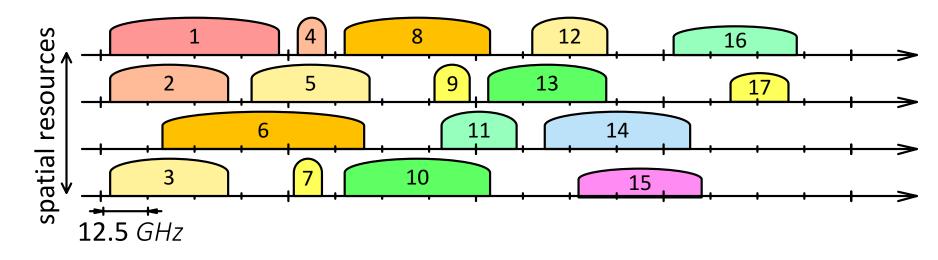
#### **SDM Scenarios**

[Spectral dimension/spatial dimension]

- Flexgrid/Single parallel transmission in EON network
- Flexgrid/Fixed SSChs can be transmitted using different SpRcs, however, within the same spectrum segment
- **Flexgrid/Flexible** full spectral and spatial flexibility in forming SSChs is allowed. Although this scenario enables best resource utilization, it may lead to fragmentation of spectrum



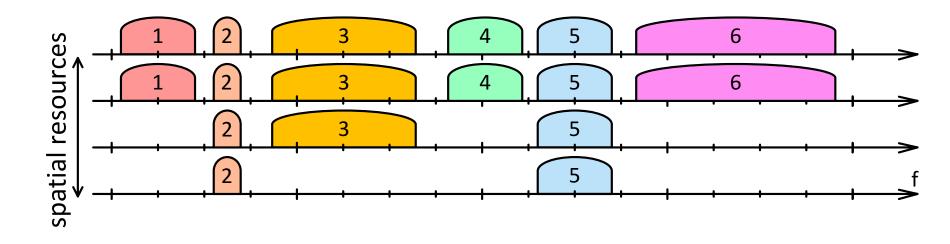
#### **Flexgrid/Single Scenario**



**Independent switching** of spatial mode and wavelength channel (i.e., space–wavelength granularity)



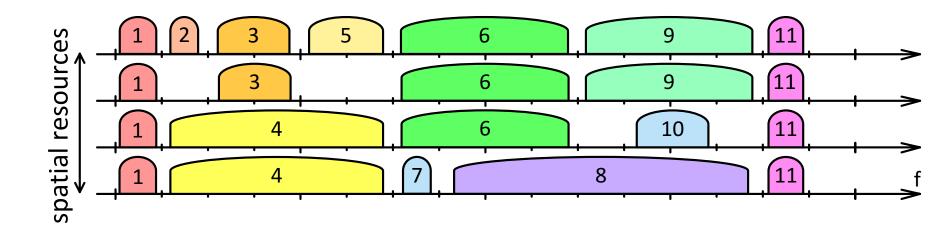
#### **Flex-grid/Fixed Scenario**



Wavelength switching across all spatial modes (wavelength granularity), also called spectral switching or joint switching



#### **Flex-grid/Flexible Scenario**



Independent switching or wavelength switching across spatial mode subgroups (fractional space– full wavelength granularity), sometimes called fractional joint switching or grouped spectral switching



#### **Pros and Cons of SDM**

- Increase the overall transmission capacity of optical networks beyond the limits of WDM and EON networks in a cost-effective manner by integrating the SDM equipment (transceivers, switching devices) to enable to a certain extent realizing multiple transmission systems in parallel
- ③ All advances of EONs can be used in SDM networks
- New fibers are required for multi-core or multi-mode transmissions
- Key network components for SDM (amplifiers, multiplexers, transceivers) are under development
- <sup>(2)</sup> Crosstalks between cores/modes can limit transmission range



### **Goal and Novelty**

- The main **goals** of this work are :
  - To develop an effective heuristic algorithm for the Flexible scenario
  - To examine main characteristics of the Flexible scenario in terms of the spectrum usage



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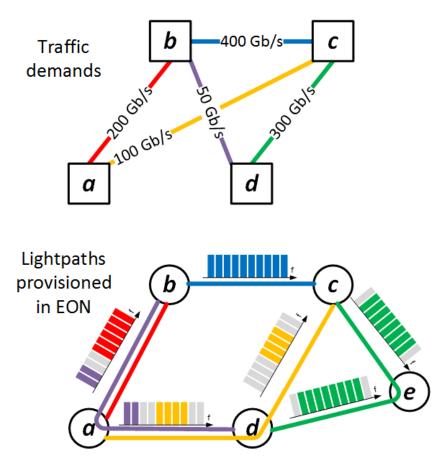
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#### **Routing and Spectrum Allocation (RSA)**

The basic optimization problem in EONs is RSA (Routing and Spectrum Allocation) that consists in selection for every demand of a routing path and spectrum with the following constraints:

- Continuity constraint states that in an absence of spectrum converters, the demand must use exactly the same spectrum slots (optical corridor) in all links included in the routing path
- Contiguity constraint requires that slices assigned to a particular demand must be adjacent (contiguous)





#### **Distance-Adaptive Transmission**

- In SDM networks based on the concept of EONs, it is possible to use various modulation formats, e.g., BPSK, QPSK, 8-QAM, 16-QAM
- These modulation formats provide some trade-off between spectrum efficiency and transmission range, i.e., more spectrum effective modulation formats provide shorter transmission range
- A reasonable approach is a distance-adaptive transmission (DAT), i.e., a modulation format for a particular demand is preselected based only on the transmission distance



#### **DAT - Example**

Distance-Adaptive Modulation Formats for Bit-Rate 400 Gb/s

|               | BPSK  | QPSK  | 8-QAM | 16-QAM |
|---------------|-------|-------|-------|--------|
| #transceivers | 8     | 4     | 3     | 2      |
| #slices       | 25    | 13    | 10    | 7      |
| Range [km]    | 6 300 | 3 500 | 1 200 | 600    |

Path length 900 km -> 8-QAM, 3 transceivers, 10 slices Path length 1800 km -> QPSK, 4 transceivers, 13 slices



#### **Demand Provisioning in SDM**

- Demand on the selected path is assigned to a *spectral-spatial channel* (SSCh) using spectral resources that can be allocated on more than one SpRc
- In consequence, the **number of possible (SSChs**) in SDM networks is much **larger** comparing channels in **EONs**
- The basic optimization problem in SDM networks is RSSA (Routing, Space and Spectrum Allocation)

[Walkowiak K., Lechowicz P., Klinkowski M., Sen A., ILP Modeling of Flexgrid SDM Optical Networks, Proceedings of the 17<sup>th</sup> International Telecommunications Network Strategy and Planning Symposium **Networks 2016**, (Montreal, Canada, September 26-28, 2016), 2016]



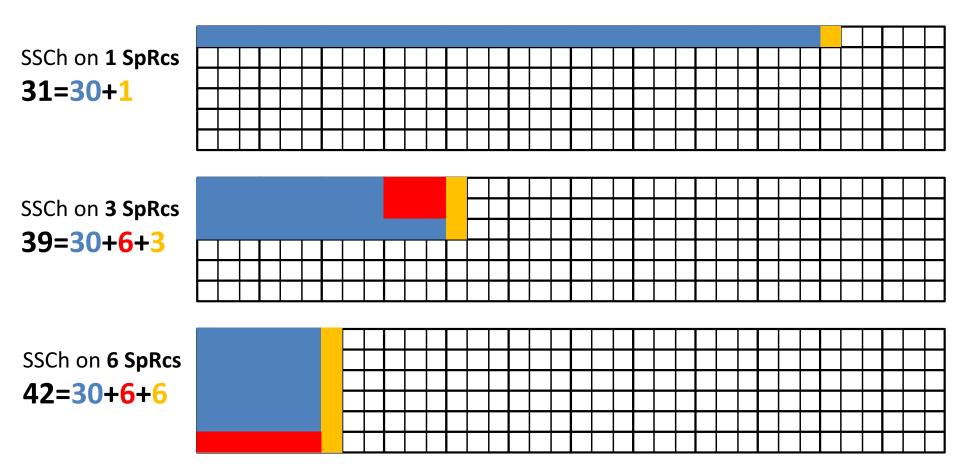
# SDM – Example (1)

- Demand bit-rate is **1 Tbps**
- Path length is 3000 km
- According to DAT, the selected MF is **QPSK**
- Since QPSK supports 100 Gbps per one transceiver, we need
  10 transceivers (=1 Tbps/100 Gbps) and 30 slices of 12.5 GHz



#### SDM – Example (2)

30 slices required to etablish 1 Tbps demand using QPSK on 3000 km path





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#### ILP Model

#### objective

min  $\sum_{s \in S} y_s$ 

#### constraints

$$\begin{split} \sum_{p \in P(d)} \sum_{c \in C(d,p)} x_{dpc} &= 1 & d \in D \\ \sum_{d \in D} \sum_{p \in P(d)} \sum_{c \in C(d,p)} \gamma_{dpcsk} \delta_{edp} x_{dpc} &\leq y_{esk} & e \in E, \ k \in K(e), \ s \in S \\ \sum_{k \in K(e)} y_{esk} &\leq |K(e)| y_{es} & e \in E, \ s \in S \\ \sum_{e \in E} y_{es} &\leq |E| y_s & s \in S \end{split}$$



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# **Greedy Algorithm**

**Require:** set of demands *D*, sets *P*(*d*) with candidate paths for each demand , SSCh comparing strategy *comp*, sorting type *sort* 

1 **function** *Greedy*(*D*, *P*, *comp*, *sort*)

- 2 D := sortDemands(D, sort)
- 3 **for** *i* := 0 **to** |*D*| **do**
- $4 \qquad d:=D[i]$
- 5 [p, ssch] := FPCSpectrum(P(d); comp)
- 6 allocate(*p*, *ssch*)



## **Tuning - Sorting**

As sort, we consider one of the following metrics:

- **Slices** the required number of slices on the shortest path
- **Distance** the length (in km) of the demand's shortest path
- **Hop count** the number of links on the shortest path



# **Tuning – SSCh Selection**

- Lowest Start (LS) the SSCh of the lowest starting slice
- Lowest End (LE) the SSCh of the lowest ending slice index is selected
- Penalty (PEN) the SSCh with the lowest penalty  $\Theta_1(c)$  is selected:  $\Theta_1(c) = \alpha (guardband(c) + rounding(c)) + end(c)$
- Demands-Varying Penalty (DVP) the SSCh with the lowest penalty  $\Theta_2(c)$  is selected:

 $\Theta_2(c) = \alpha \cdot (1 - \tau) \cdot (guardband(c) + rounding(c)) + (1 - \alpha) \cdot \tau \cdot end(c)$ 

- end(SSCh) returns an index of the highest slice used by SSCh
- rounding(SSCh) returns the amount of slices wasted for rounding
- guardband(SSCh) returns the amount of slices used for guardbands
- au is equal to the ratio of currently allocated demands to all demands



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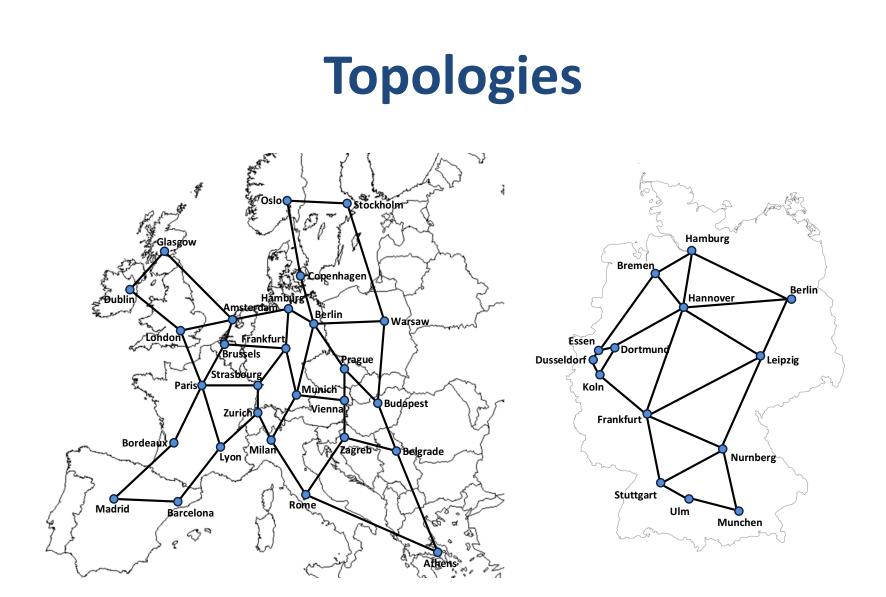
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#### Assumptions

- Transceivers operate at fixed baud rate of 28 GBaud and each transceiver transmits/receives an optical channel (optical carrier) that occupies 3 slices of 12.5 GHz
- A fixed **guardband** defined as 1 slice of 12.5 GHz
- Four modulation formats: BPSK, QPSK, 8-QAM, and 16-QAM with range 6300 km, 3500 km, 1200 km and 600 km, with bit-rate: 50 Gbps, 100 Gbps, 150 Gbps and 200 Gbps, respectively
- Each **demand** has the bit-rate selected at random from range 50 Gbps to 1 Tbps with 50 Gbps granularity
- Number of **candidate paths** for each demand is 30







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#### **Tuning – Number of Slices**

|            | Sorting |          |           |  |
|------------|---------|----------|-----------|--|
| Algorithm  | Slice   | Distance | Hop Count |  |
| LS         | 1066.1  | 1125.8   | 1079.5    |  |
| LE         | 1300.8  | 1367.3   | 1331.3    |  |
| PEN(α=0.2) | 1300.1  | 1366.7   | 1335.5    |  |
| PEN(α=0.5) | 1301.4  | 1365.3   | 1332.6    |  |
| PEN(α=0.8) | 1264.1  | 1323     | 1288.1    |  |
| DVP(α=0.2) | 1236.7  | 1319     | 1278.6    |  |
| DVP(α=0.5) | 1202.3  | 1274.1   | 1230.6    |  |
| DVP(α=0.8) | 1159    | 1217.5   | 1157.7    |  |



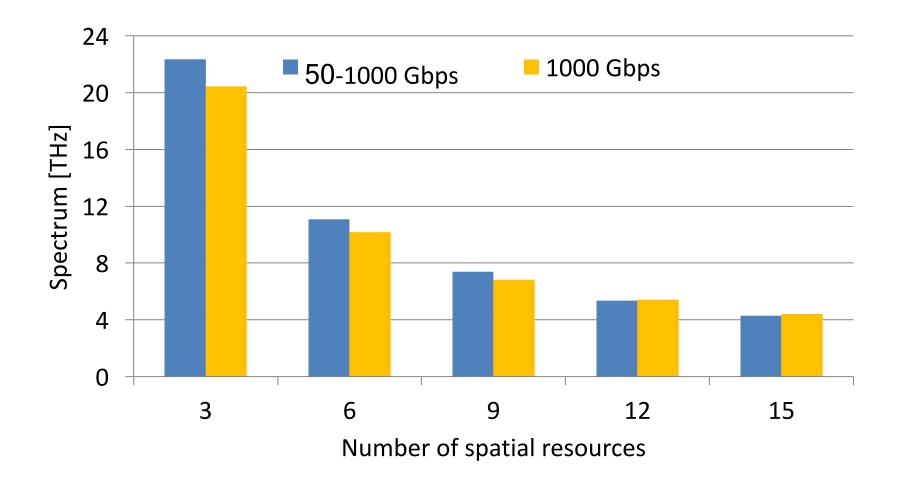
#### **CPLEX vs. Heuristic for Euro28**

|              |    | Number of slices |        | Execution time |        |
|--------------|----|------------------|--------|----------------|--------|
| <i>P</i> (d) | D  | CPLEX            | Greedy | CPLEX          | Greedy |
| 4            | 20 | 28               | 28     | 260s           | <1ms   |
| 4            | 30 | 31               | 31     | 1h             | <1ms   |
| 4            | 40 | 34               | 34     | 1h             | <1ms   |
| 4            | 50 | Out of<br>Memory | 58     | -              | <1ms   |
| 2            | 20 | 28               | 28     | 60s            | <1ms   |
| 2            | 30 | 25               | 31     | 1h             | <1ms   |
| 2            | 40 | 34               | 34     | 1h             | <1ms   |
| 2            | 50 | Out of<br>Memory | 58     | -              | <1ms   |



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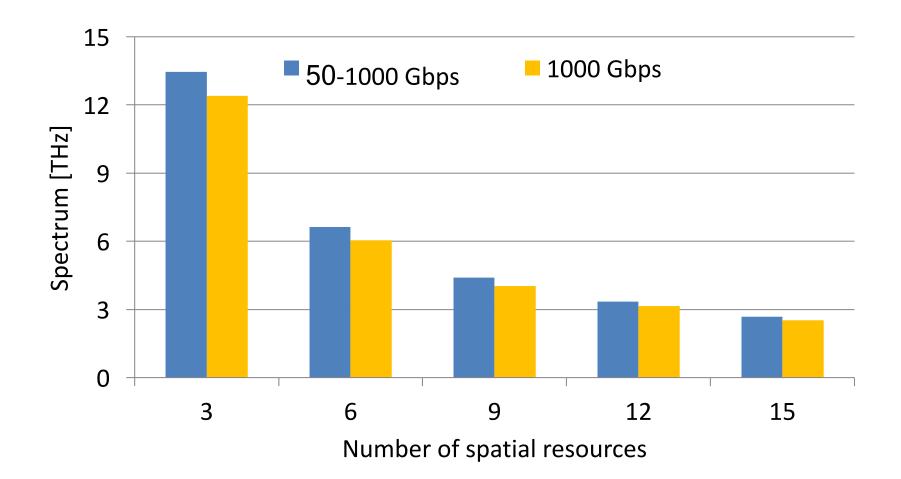
# Spectrum usage for various types of demands – network Euro28 and 1 Pbps traffic





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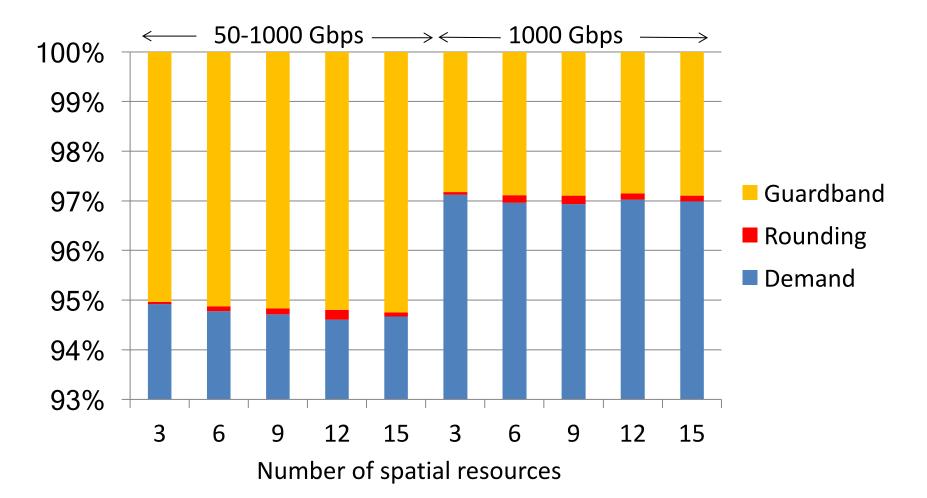
# Spectrum usage for various types of demands – network DT14 and 1 Pbps traffic





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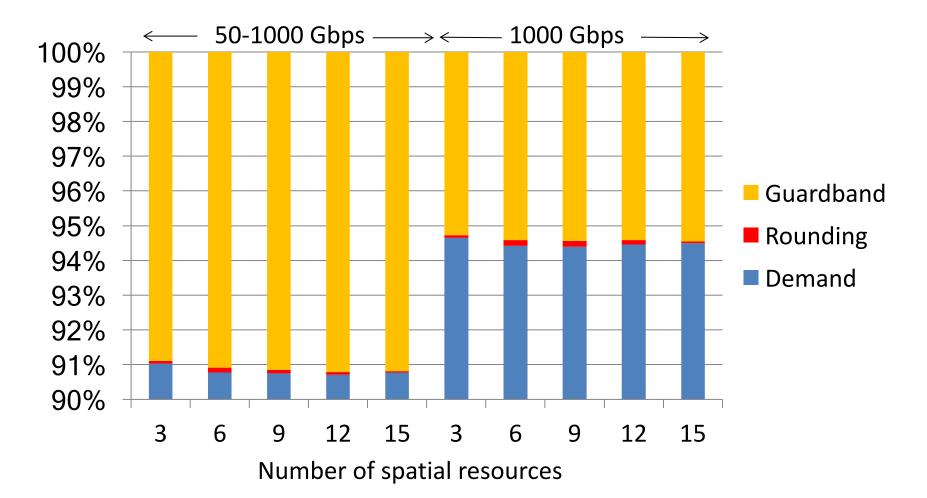
# Spectrum usage for various types of demands – network Euro28 and 1 Pbps traffic





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# Spectrum usage for various types of demands – network DT14 and 1 Pbps traffic





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#### Average execution time of the heuristic (in seconds) as a function of the number of SpRcs for, 1 Pbps traffic

|         | Number of spatial resources |    |    |     |      |
|---------|-----------------------------|----|----|-----|------|
| Network | 3                           | 6  | 9  | 12  | 15   |
| Euro28  | 2                           | 11 | 63 | 325 | 1691 |
| DT14    | 1                           | 7  | 28 | 133 | 621  |



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#### Conclusions

- A greedy algorithm with different strategies for sorting of demands and allocation of spectral-spatial channels provides results close to optimal ones
- Spectrum usage in examined topologies decrease almost proportionally with the increase of SpRcs
- The Flexible scenario yields similar results the Single scenario.
  Flexible, despite its capability to form SSChs in both domains, most of the time selects SSChs which use only one SpRc
- The **Flexible** scenario is **very complex** in terms of SSChs number, what results in **high execution time** of simple heuristic
- Future work includes development of **heuristic and metaheuristic** methods that enable solving **large problem instances**



# Thank you for attention

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