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# Bit Index Explicit Replication (BIER) Multicasting in Transport Networks

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## ❑ Segment Routing (SR)

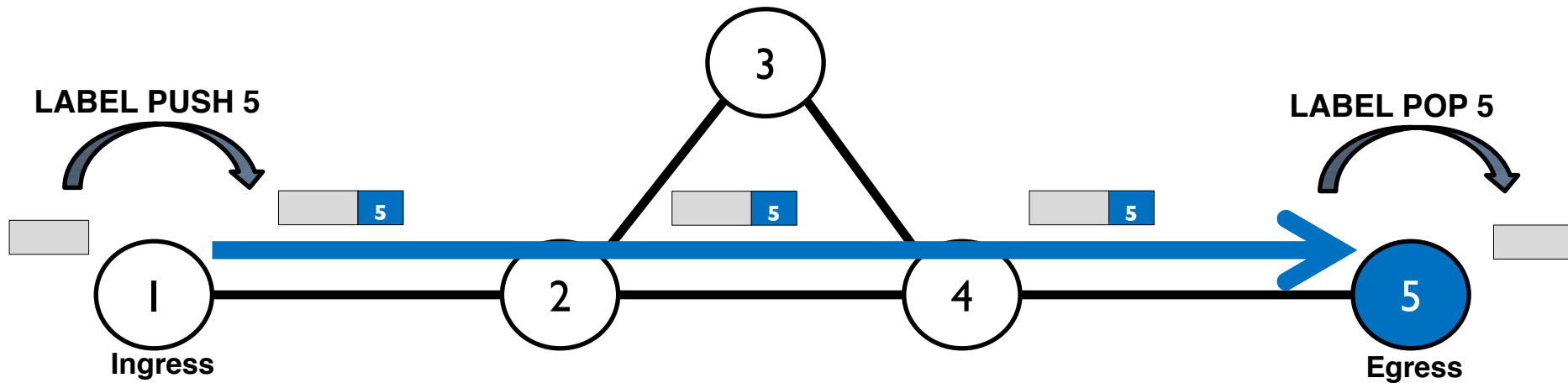
- technology
- Path encoding
- TE performance (ECMP vs strict routing)
- Use cases

## ❑ Bit Index Explicit Replication (BIER)

- technology
- Experimental validation

- ❑ **Segment Routing (SR)** is a traffic engineering (TE) technique compatible with traditional **MPLS data plane**.
- ❑ Using SR, a **signaling protocol is not required** and path state is not maintained in intermediate nodes → simplified control plane operation.
- ❑ Each packet is forwarded according to an header composed of **segment identifiers** (SIDs), e.g., representing a specific network node.
- ❑ SIDs are advertised by properly extended IGP (e.g., OSPF-TE).
- ❑ Intermediate nodes forward the received packet along the **shortest path** toward the node indicated in the top SID.

# Segment Routing (SR) basic behavior (1/2)

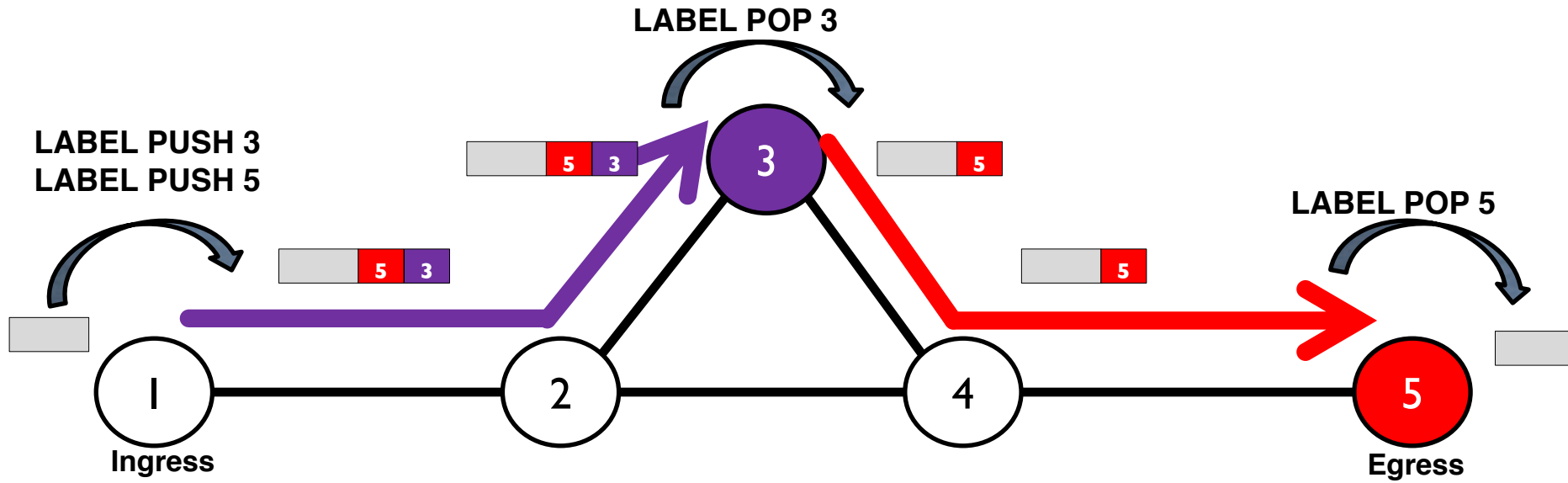


Target path: 1,2,4,5



segment list: 5

# Segment Routing (SR) basic behavior (2/2)

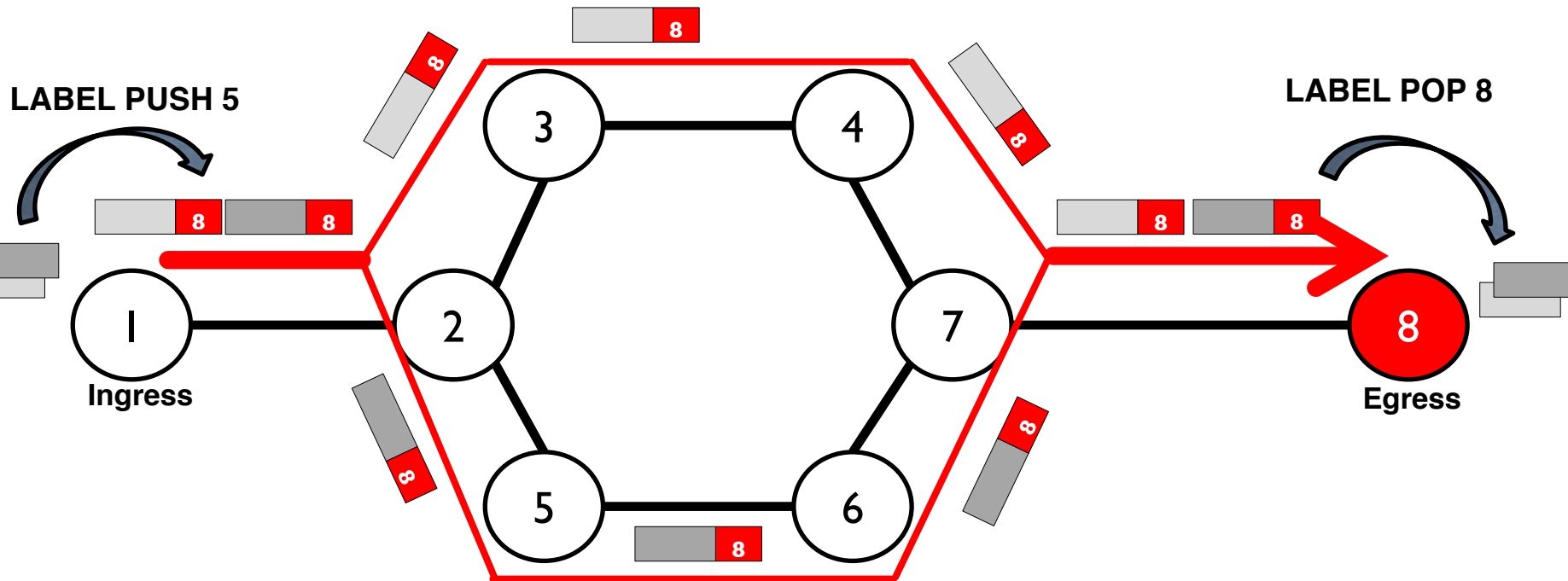


Target path: 1,2,3,4,5



segment list: 3, 5

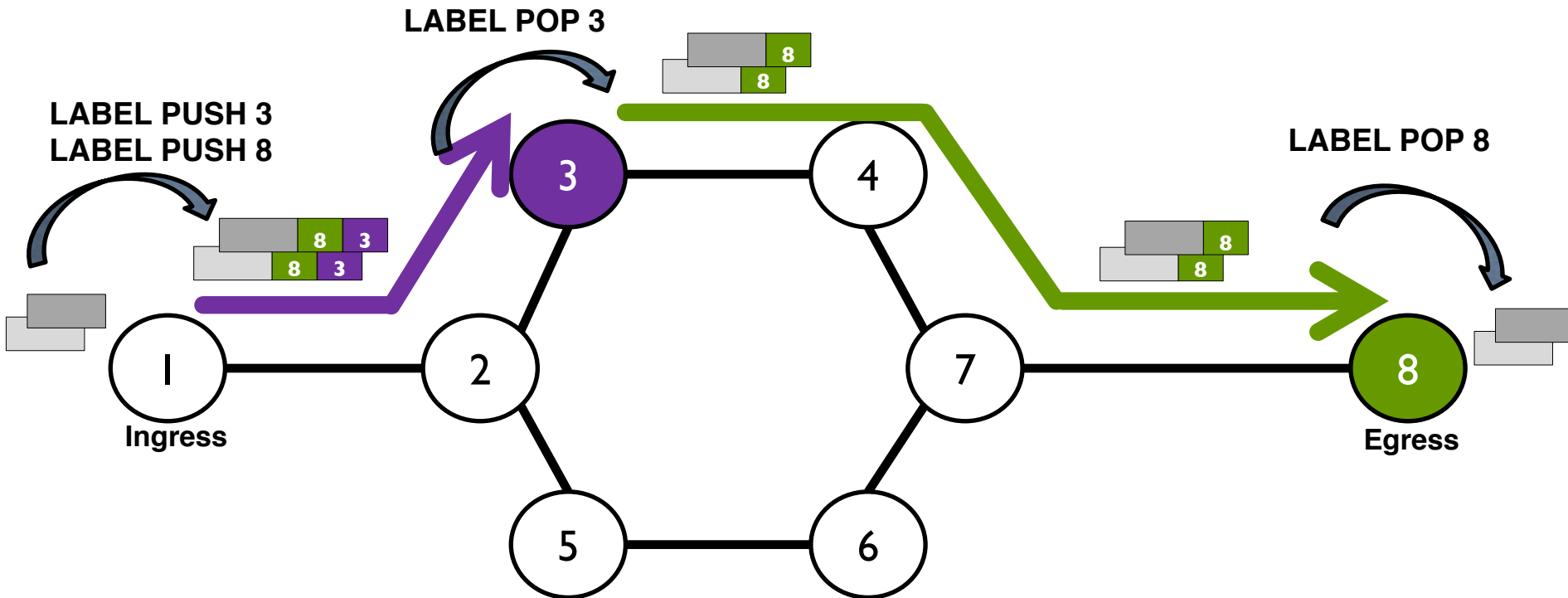
# Equal Cost Multi Paths (ECMP)



Default behavior → load balancing on ECMPs

Target paths: 1,2,3,4,5 AND 1,2,6,4,5 → segment list: 8

# Strict route selection avoiding ECMPs

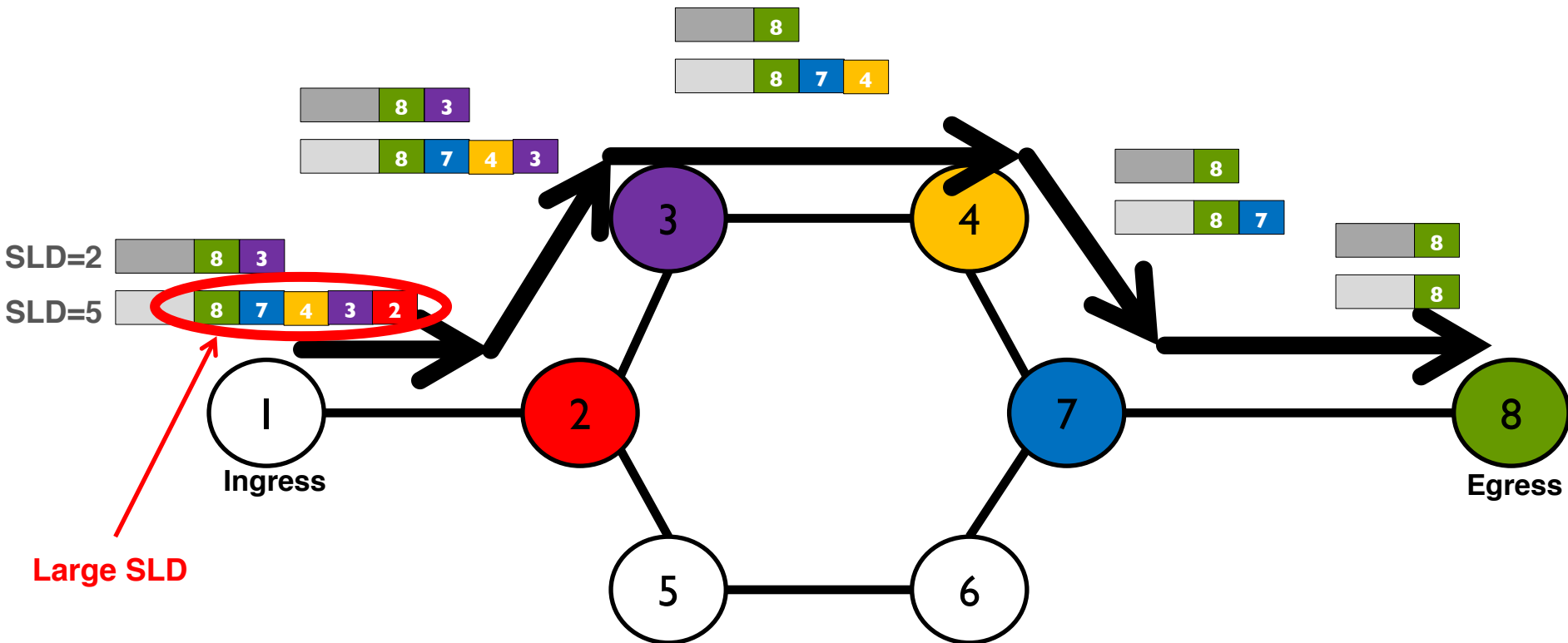


Target path: 1,2,3,4,5,7,8



segment list: 3, 8

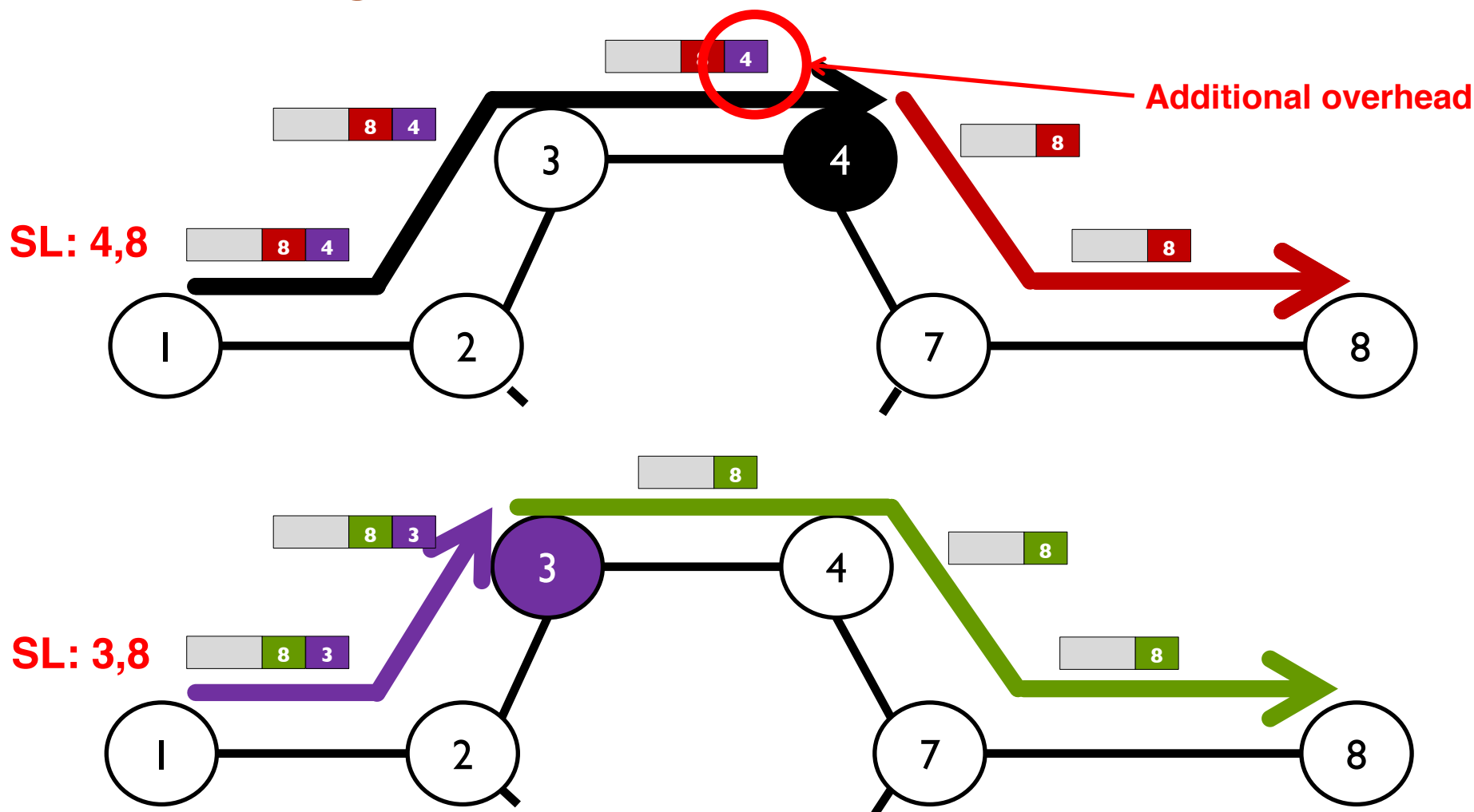
# Path encoding: Segment List with minimum SLD



- **Problem:** given a strict route, identify the segment list having the minimum Segment List Depth (SLD).
- Routers may **not** support large SLD

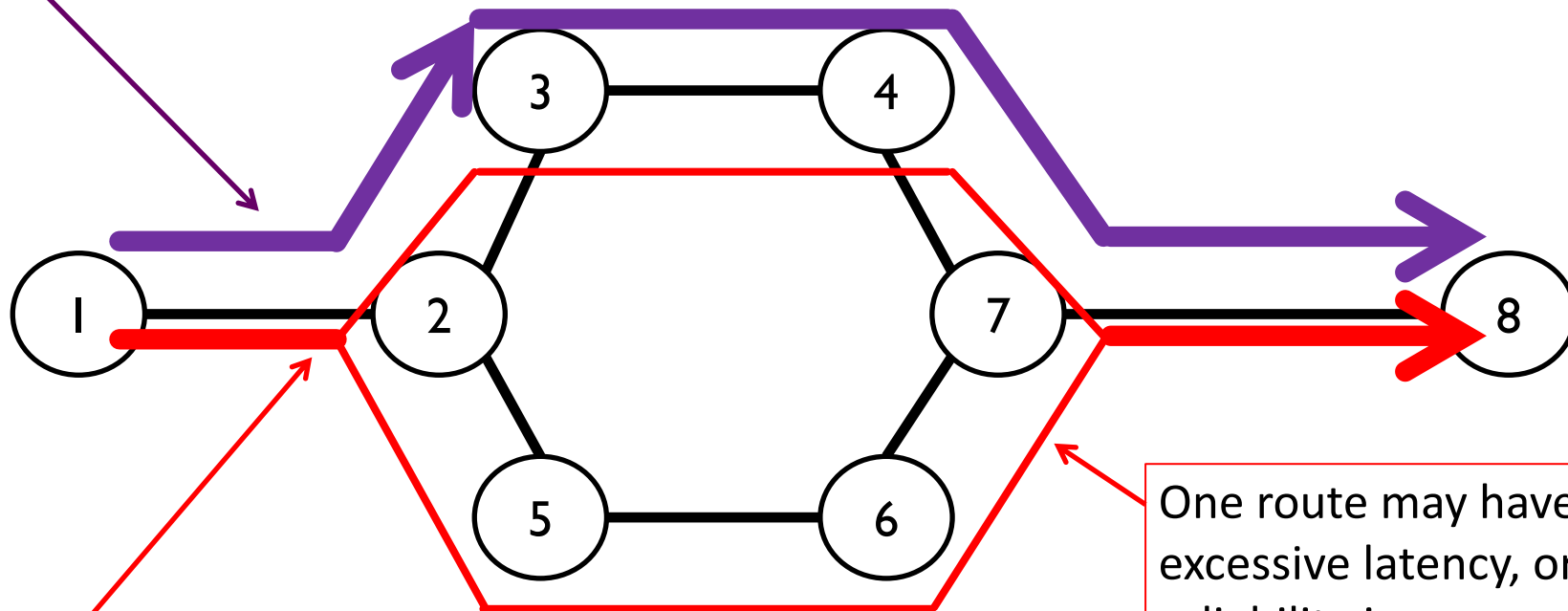


# Path encoding: Segment List with minimum overhead



# ECMP vs. Strict route

Larger segment lists



In ECMP, Deep Packet Inspection (DPI) is needed to perform load balancing on a per-flow basis, avoiding packet misordering  
→ additional HW requirements, not needed in case of strict routes

- Network usage for different topologies
- Default SR behavior exploiting ECMP typically leads to inefficient network usage
- Segment Routing with large Segment List Depth (e.g., 8) guarantees optimal TE solutions
- Surprisingly, Segment Routing with just a Segment List Depth of 3 is able to guarantee optimal TE solutions

	ECMP	SHP	SEGMR $\kappa = 8$	SEGMR $\kappa = 3$
Grid $2 \times 2$	2.00	2.00	2.00	2.00
Grid $3 \times 3$	7.00	6.00	6.00	6.00
Grid $4 \times 4$	18.63	16.00	16.00	16.00
Grid $5 \times 5$	36.75	30.00	30.00	30.00
Eurocore	4.96	4.00	4.00	4.00
NFSNET	15.33	13.00	13.00	13.00
EON	25.00	18.00	18.00	18.00
UKNET	27.63	21.00	19.00	19.00
ITALNET	48.38	33.00	28.00	28.00
Arpanet	35.88	33.00	33.00	33.00
Eurolarge	131.60	88.04	66.00	66.00

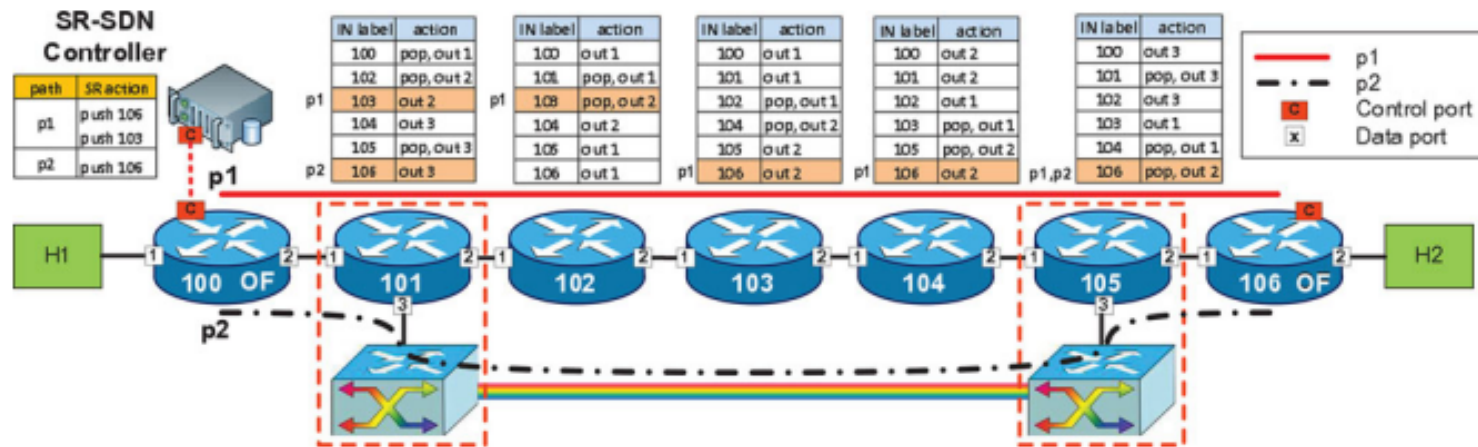
Number of flows in the most utilized link

- ECMP: default SR behavior ( $k=1$ )
- SHP: shortest path, no ECMP
- SEGMR: SR with  $k=8$  and  $k=3$

$k$ : segment list depth

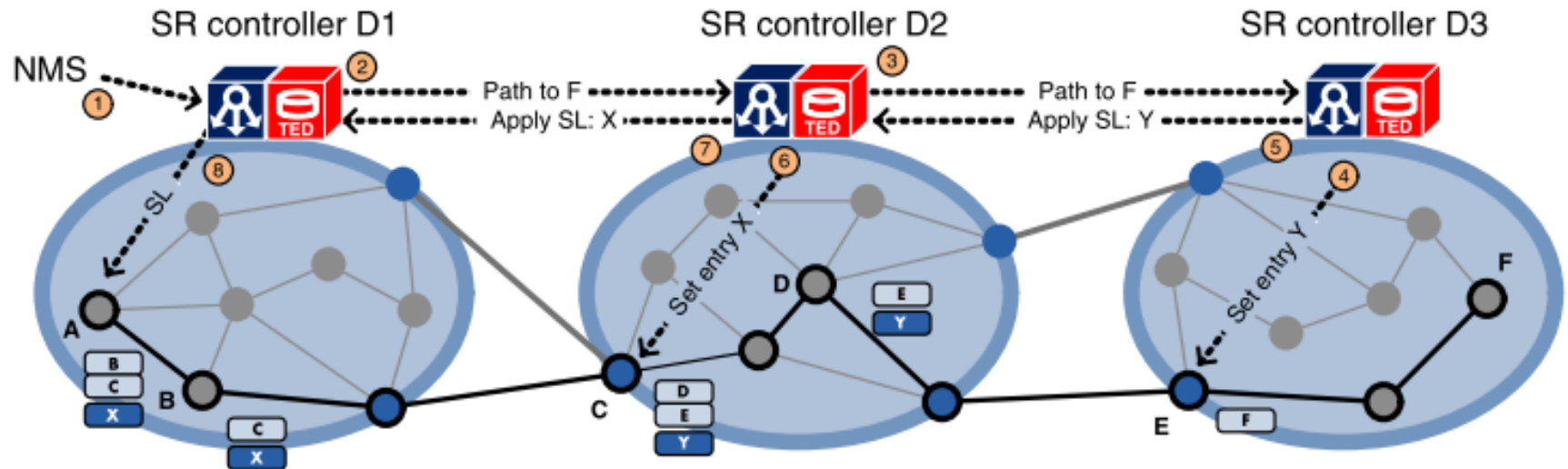
Uniform traffic matrix is assumed. Results are confirmed for non-uniform matrices

# Use cases for SR: Optical bypass



- Dynamic selection of pre-established optical bypass, with no signaling

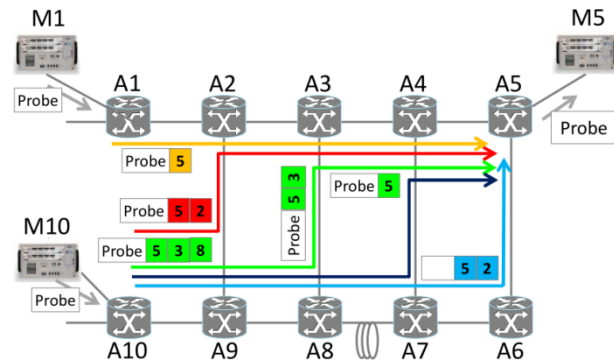
# Use cases for SR: multi-domain



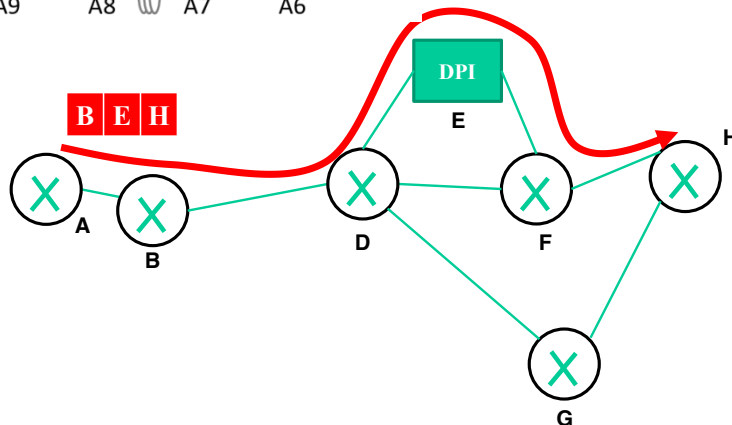
- No e2e signaling (critical in multi-domain/vendor scenarios)
- Compressed Segment Lists to limit label stacking
- Confidentiality

- Recovery:
  - SR-FAILOVER: rerouting to the destination
  - SR-DETOUR: rerouting to the next(-next) hop

- OAM



- Service chaining



# Bit Index Explicit Replication (BIER)

- ❑ BIER has been recently proposed for P2MP
- ❑ As in SR,
  - No signaling protocol
  - No forwarding state at intermediated nodes
  - the ingress router applies a specifically designed label (here called BitString) which defines the forwarding actions
- ❑ In the BIER BitString, each bit represents exactly one egress router in the domain.
- ❑ Forwarding is then performed by each intermediate node by just processing and updating the BitString
- ❑ In large networks, a hierarchical structure of the BIER header is used.

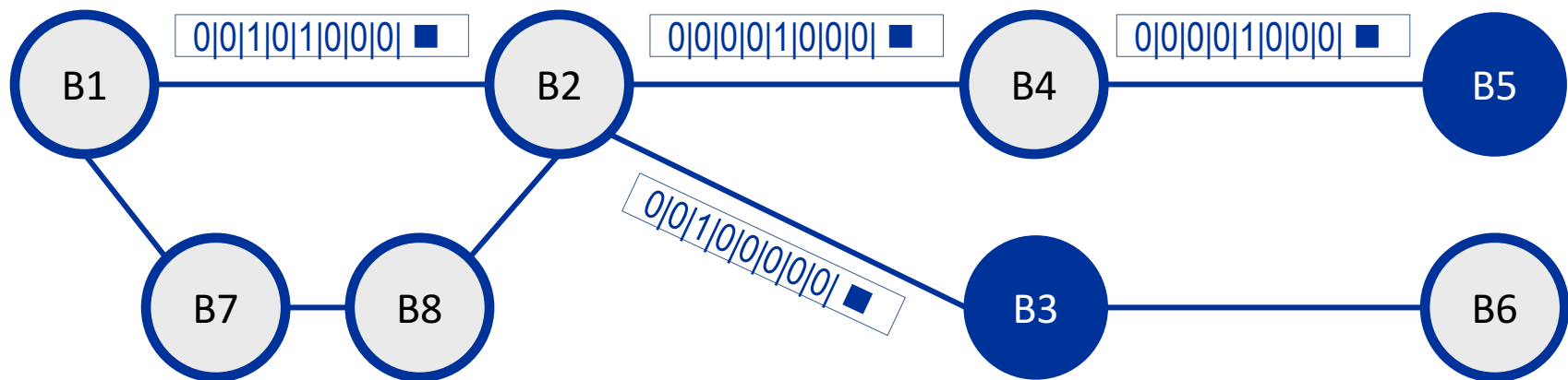
# BIER

## Basic behavior

BIT INDEXED FORWARDING TABLES

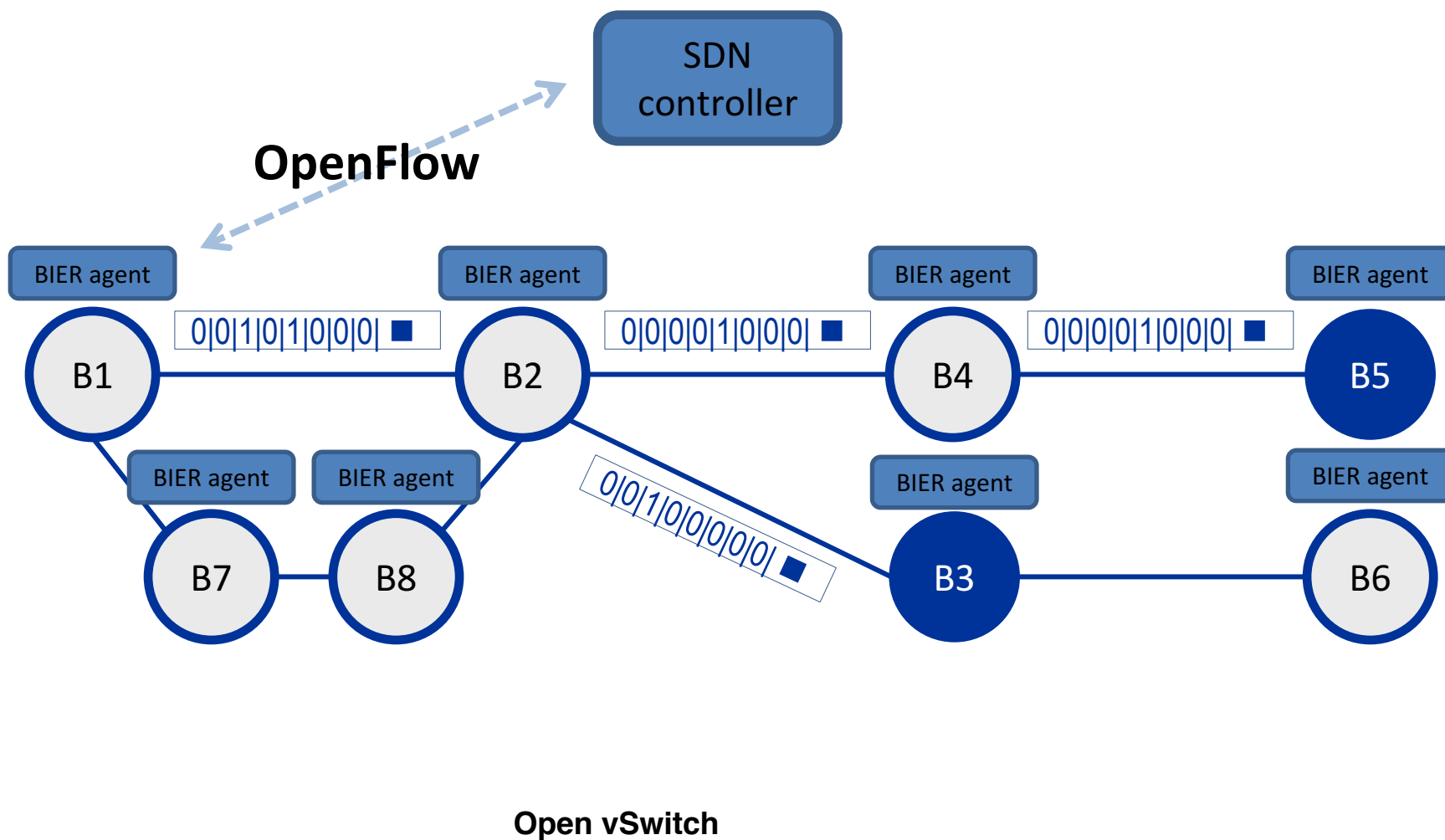
BIFT at <i>B2</i>			BIFT at <i>B3</i>		
BitMask	Neighbour	Port	BitMask	Neighbour	Port
00000001	<i>B1</i>	1	00000100	-	local
00000010	-	local	11011011	<i>B2</i>	1
00011000	<i>B4</i>	2	00100000	<i>B6</i>	2
00100100	<i>B3</i>	3	-	-	-
11000000	<i>B8</i>	4	-	-	-

BIFT table scales with the number of outgoing links and not with the number of flows traversing the node.





## Implementation with OF controller



# BIER implementation

- SDN network controller implemented in Ryu, with OpenFlow 1.3
- Packet nodes: Open vSwitch + a BIER agent in each node implemented on a local instance of the Ryu controller
- The BIER agent stores the list of nodes reachable with a shortest path using each of the outgoing links (BIFT table - provided and updated by the network SDN controller.
- When a multicast request arrives, the SDN controller only has to enforce the proper BIER header at the ingress (no signaling protocol).
- A group is created with the *id* equal to the numerical value of the BitString including all recipients joining the multicast address.
- If packets need to be replicated at the specific node, the group will include a number of buckets.

Node 2: Flow Table		Node 2: Group Table	
Flow match	Action	Group #	Action
MPLS label 10100	group 10100	10100	Bucket 1: set label 00100, output port 3
			Bucket 2: set label 10000, output port 2
MPLS label 10010	group 10010	10010	Bucket 1: pop label, port local
			Bucket 2: set label 10000, output port 2

# Capture

capture\_ondm.pcap

File Modifica Visualizza Vai Cattura Analizza Statistiche Telefonia Wireless Strumenti Aiuto

openflow\_v4

Espressione...

No.	Time	Source	Destination	Protocol	Info
1	0.000	10.30.2.36	10.30.2.36	OpenFlow	Type: OFPT_ECHO_REQUEST
2	0.000	10.30.2.36	10.30.2.36	OpenFlow	Type: OFPT_ECHO_REPLY
4	4.223	10.30.2.36	10.30.2.36	OpenFlow	Type: OFPT_PACKET_IN
5	4.227	10.30.2.36	10.30.2.36	OpenFlow	Type: OFPT_GROUP_MOD
7	4.227	10.30.2.36	10.30.2.36	OpenFlow	Type: OFPT_FLOW_MOD
9	4.227	10.30.2.36	10.30.2.36	OpenFlow	Type: OFPT_PACKET_OUT

(a) - message exchange

OpenFlow 1.3

Version: 1.3 (0x04)  
Type: OFPT\_GROUP\_MOD (15)  
Length: 112  
Transaction ID: 282450948  
Command: OFPGC\_ADD (0)  
Type: OFPGT\_ALL (0)  
Pad: 00

**GROUP\_MOD**  
**Node C**

Group ID: 12

Bucket

Length: 48  
Weight: 0  
Match port: OFPP\_ANY (0xffffffff)  
Match group: OFPG\_ANY (0xffffffff)  
Pad: 00000000

Action

Type: OFPAT\_SET\_FIELD (25)  
Length: 16

OXM field

Class: OFPXNC\_OPENFLOW\_BASIC (0x8000)  
0100 010. = Field: OFPXMT\_OFB\_MPLS\_LABEL (34)  
.... ..0 = Has mask: False  
Length: 4  
Value: 8

Pad: 00000000

Action

Bucket

Out port: 0  
Out group: 0  
Flags: 0x0000  
Pad: 0000

Match

Type: OFPMT\_OXM (1)  
Length: 18

OXM field

Class: OFPXNC\_OPENFLOW\_BASIC (0x8000)  
0000 101. = Field: OFPXMT\_OFB\_ETH\_TYPE (5)  
.... ..0 = Has mask: False  
Length: 2  
Value: MPLS label switched packet (0x8847)

OXM field

Class: OFPXNC\_OPENFLOW\_BASIC (0x8000)  
0100 010. = Field: OFPXMT\_OFB\_MPLS\_LABEL (34)  
.... ..0 = Has mask: False  
Length: 4  
Value: 12

Pad: 000000000000

Instruction

Type: OFPIT\_APPLY\_ACTIONS (4)  
Length: 16  
Pad: 00000000

Action

Type: OFPAT\_GROUP (22)  
Length: 8  
Group ID: 12

(b)

(c)

FLOW\_MOD  
Node C

Pacchetti: 19 - visualizzati: 11 (57.9%) - marcati: 2 (10.5%) - Tempo di caricamento: Accesso a Internet

# Conclusions

- Overview of Segment Routing technology and use cases
  - Path encoding
  - traffic engineering (limited SLD are typically adequate)
  - Multi-domain (no e2e signaling)
  - Recovery (straightforward)
  - OAM (easy mechanism to probe the network)
  - Service chaining (services can be described with Segment ID)
- BIER
  - Experimental demonstration

# thank you!

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